http: / /www-ai.cs.uni-dortmund.de/weblab/static/RLP/html/


RELOOP: A Toolkit for Relational Convex Optimization


Artificial Intelligence and Machine Learring



David
Poole

## MORGAN\&CLAYPOOL PUBLISHERS

Statistical Relational
Artificial Intelligence

and many more ...

# The Democratization of Optimization 

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Kristian Kersting

## Democratization of Data



## Take your spreadsheet ...

Features

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## IS IT REALLY THAT SIMPLE?

## NO, e.g., today's data is relational




Relation 1


Relation 2

Objects


Relation 2
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Meart diseases and strokes cardiovascular disease - are expensive for the world

According to the World Heart Federation, cardiovascular disease cost the European Union EURO169 billion in 2003 and the USA about EUR0310.23 billion in direct and indirect annual costs. By comparison, the estimated cost of all cancers is EURO1.46.19 billion and HIV infections, EURO22.24 billion
nature
REVIEWS
nature
REVIEWS
 Electronic Health Records A New Oppoftunity for AI to Save our Lifes
[Natarajoan, Kersting, et al. IAAI 2013, Springer Briefs in CS 2015, AIME 2015]

## EHRs are dirty and interconnected



## Statistical Relational AI

... the study and design of intelligent agents that act in noisy worlds composed of objects and relations among the objects


Thanks to you - the Iatlian AI community - for your great contributions!

## Boosted Statistical Relational Learning

Atherosclerosis is the cause of the majority of Acute Myocardial Infarctions (heart attacks)

[Circulation; 92(8), 2157-62, 1995; JACC; 43, 842-7, 2004]

| Algorithm | Accuracy | AUC-ROC |
| :---: | :---: | :---: |
| J48 | 0.667 | 0.607 |
| SVM | 0.667 | 0.5 |
| AdaBoost | 0.667 | 0.608 |
| Bagging | 0.677 | 0.613 |
| NB | 0.75 | 0.653 |
| RPT | $0.669^{*}$ | 0.778 |
| RFGB | $0.667^{*}$ | 0.819 |


| Algo | Likelihood | AUC-ROC | AUC-PR | Time |
| :---: | :---: | :---: | :---: | :---: |
| Boosting | 0.810 | 0.961 | 0.930 | 9s |
| MLN | 0.730 | 0.535 | 0.621 | $\mathbf{9 3}$ hrs |

## Take-Away Messages

1. Graphical models allow to deal with uncertainty and to make predictions
2. Graphs/Matrices are not enough, we need logic/high-level languages

## STILL NOT CONVINCED?

## Guy van den Broeck's not

 so simple AI example first card of a shuffled deck is an Ace?

## Easy for humans but not so easy for graphical models

Exact inference builds a table of $\geq 13^{52}$ rows!

Message passing passes $\geq 13^{52}$ messages!

Graphical model is fully connected, no independencies, high tree-width

# Low tree-width is not the final answer 

## Fast modelling

## Fast inference

We need relationa nodels
$w 1: \forall p, x, y, \operatorname{Card}(p, x) \wedge \quad \operatorname{ard}(p, y) \Rightarrow x=y$
$w 2: \forall c, x, y, \operatorname{Card}(x, c) \wedge \operatorname{Card}(y, c) \Rightarrow x=y$
and symmetry-reduction

## Take-away Messages

# 1. Graphical models allow to deal with uncertainty 

2. Graphs/Matrices are not enough, we need logic /high-level languages
3. Tree-Width is not the end of the story

[Singla, Domingos AAAI 2008; Kersting, Ahmadi, Natarajan UAI 2009; Ahmadi, Kersting, Mladenov, Natarajan ML] 2013]

## Lifted Loopy Belief Propagation = Exploit computational symmetries



[^0]
## Compression: Coloring the graph

- Color nodes according to the evidence you have

- No evidence, say red
- State „one", say brown
- State „two", say orange
" ...
- Color factors distinctively according to their equivalence classes. For instance, assuming $f_{1}$ and $f_{2}$ to be identical and $B$ appears at the second position within both, say blue


## Compression: Pass the colors around



1. Each factor collects the colors of its neighboring nodes
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4. Nodes are recolored according to the collected signatures


## Compression: Pass the colors around



1. Each factor collects the colors of its neighboring nodes
2. Each factor "signs" ist color signature with its own color
3. Each node collects the signatures of its neighboring factors
4. Nodes are recolored according to the collected signatures
5. If no new color is created stop, otherwise go back to 1

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## Compression: <br> ... and compute the quotient factor graph



## Essentially we just compute the so-called quotient factor graph

## Finally, run a modified Loopy Belief Propagation



- Nodes are now groups of random variables
- The counts ensure that we send the same number of message as standard loopy belief propagation


## Lifted Probabilistic Inference

## Compress the model

## Run a modified inference method

might also be interwined


## Weisfeiler-Lehman (WL) Algorithmus aka "naive vertex classification"

- Basic subroutine for GI testing
- Computes LP-relaxations of GA-ILP, aka. fractional automorphisms
- Quasi-linear running time $\mathrm{O}((\mathrm{n}+\mathrm{m}) \log (\mathrm{n}))$ when using
 asynchronous updates [Berkholz, Bonsma, Grohe ESA 2013]
- Part of graph tool SAUCY [See e.g. Darga, Sakallah, Markov DAC 2008]
- Can be extended to weighted graphs/real-valued matrices [Grohe, Kersting, Madenov, Selman ESA 2014]
- Actually a Frank-Wolfe optimizer and can be viewed as recursive spectral clustering [Kersting, Mladenov, Garnett, Grohe AAAI 2014]


## From Factor Graphs to Graphs



Encode the factor colors into the node colors


Then run Weisfeiler Lehman / Color Passing just on the graph

[Mladenov, Ahmadi, Kersting AISTATS 2012, Grohe, Kersting, Mladenov, Selman ESA 2014, Kersting Mladenov, Tokmatov AIJ 2015]

## Compressing Linear Programs

$$
\max _{[x, y, z]^{T} \in \mathbb{R}^{3}} \quad 0 x+0 y+1 z
$$

s.t.

$$
\left[\begin{array}{ccc}
1 & 1 & 1 \\
-1 & 0 & 0 \\
0 & -1 & 0 \\
1 & 1 & -1
\end{array}\right]\left[\begin{array}{l}
x \\
y \\
z
\end{array}\right] \leq\left[\begin{array}{c}
1 \\
0 \\
0 \\
-1
\end{array}\right]
$$



## (1) Reduce the LP by running WL on the LP-Graph (2) Run any solver on the (hopefully) smaller LP



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quasi-linear overhead that may result in exponential speed up


## As also noted by Stephen Boyd

## DENSE VS. SPARSE IS NOT ENOUGH, SOLVERS NEED TO BE AWARE OF SYMMETRIES


(a)

(b)

(c)

WWIII
Feasible region of LP and the objective vectors
fractional automorpishm of the LP

Projections of the feasible region onto the span of the fractional automorphism
Compute Equitable Partition (EO) of the LP using WL


$$
\mathcal{P}=\frac{\left\{P_{1}, \ldots, P_{p}\right.}{\begin{array}{c}
\text { Partition of } \\
\text { LP variables }
\end{array}} \frac{\left.Q_{1}, \ldots, Q_{q}\right\}}{\begin{array}{l}
\text { Partition of } \\
\text { LP constraints }
\end{array}}
$$

# Intuitively, we group together variables resp. constraints that interact in the very same way in the LP. 

## Fractional Automorphisms of LPs

The EP induces a fractional automorphism of the coefficient matrix $\mathbf{A}$

$$
\mathbf{X}_{Q} \mathbf{A}=\mathbf{A} \mathbf{X}_{P}
$$

where $\mathbf{X}_{\mathrm{Q}}$ and $\mathbf{X}_{\mathrm{p}}$ are doubly-stochastic matrixes (relaxed form of automorphism)

$$
\begin{aligned}
& \left(\mathbf{X}_{P}\right)_{i j}= \begin{cases}1 /|P| & \text { if both vertices } i, j \text { are in the same } P \\
0 & \text { otherwise }\end{cases} \\
& \left(\mathbf{X}_{Q}\right)_{i j}= \begin{cases}1 /|Q| & \text { if both vertices } i, j \text { are in the same } Q \\
0 & \text { otherwise }\end{cases}
\end{aligned}
$$

## Fractional Automorphisms Preserve Solutions

If $\mathbf{x}$ is feasible, then $\mathbf{X}_{\mathrm{p}} \mathbf{x}$ is feasible, too. By induction, one can show that left-multiplying with a double-stochastic matrix preserves directions of inequalities. Hence,

$$
\mathbf{A} \mathbf{x} \leq \mathbf{b} \Rightarrow \mathbf{X}_{Q} \mathbf{A} \mathbf{x} \leq \mathbf{X}_{Q} \mathbf{b} \Leftrightarrow \mathbf{A} \mathbf{X}_{P} \mathbf{x} \leq \mathbf{b}
$$

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## Fractional Automorphisms Preserve Solutions

If $\mathbf{x}^{*}$ is optimal, then $\mathbf{X}_{\mathrm{p}} \mathbf{x}^{*}$ is optimal, too.
Since by construnction $\mathbf{c}^{T} \mathbf{X}_{P}=\mathbf{c}^{T}$ and hence $\mathbf{c}^{T}\left(\mathbf{X}_{P} \mathbf{x}\right)=\mathbf{c}^{T} \mathbf{x}$

## What have we established so far?

Instead of considering the original LP

$$
(\mathbf{A}, \mathbf{b}, \mathbf{c})
$$

It is sufficient to consider

$$
\left(\mathbf{A} \mathbf{X}_{P}, \mathbf{b}, \mathbf{X}_{P}^{T} \mathbf{c}\right)
$$

i.e. we "average" parts of the polytope.

## But why is this dimensionality reduction?

## Dimensionality Reduction

The doubly-stochastic matrix $\mathbf{X}_{P}$ can be written
as

$$
\mathbf{X}_{P}=\mathbf{B B}^{T}
$$

$$
\mathbf{B}_{i P}= \begin{cases}\frac{1}{\sqrt{|P|}} & \text { if vertex } i \text { belongs to part } P, \\ 0 & \text { otherwise. }\end{cases}
$$

Since the column space of B is equivalent to the span of $\mathbf{X}_{P}$, it is actually sufficient to consider only

$$
\left(\mathbf{A} \mathbf{B}_{P}, \mathbf{b}, \mathbf{B}_{P}^{T} \mathbf{c}\right)
$$

## This is of reduced size and actually we can also drop any constraint that becomes identical

## Dimensionality Reduction of LPs



(b)

(c)

Feasible region of LP and the objective vectors

Span of the
fractional automorpishm of the LP

Projections of the feasible region onto the span of the fractional automorphism


## Any MAP-LP message-passing approach is liftable


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## Any concave free energy is

 liftable

(a) Complete Graph MLN.


(b) Clique-Cycle MLN.


(c) Friends-smokers MLN.

## Actually, this is the first distributed lifted message-passing approach

## All MP Inference Approaches are Liftable



[Mladenov, Globerson, Kersting UAI 2014; Mladenov, Kersting UAI 2015]

## Key Idea: Refine self-loops



# coloring 

## lifting

## refine

## Lifted inference = Inference in a smaller, reparameterized model




## Classification using LPs

$$
H^{*}=\left\{\vec{x} \mid\langle\vec{x}, \vec{\beta}\rangle+\beta_{0}=0\right\}
$$



$$
d\left(H_{1}, H_{2}\right)=\frac{2}{\|\vec{\beta}\|}
$$

Replace $I_{2}$ - by $I_{1}-I_{\infty}$-norm in the standard SVM formulation

## Declarative Machine Learning

```
var pred/1;
var slack/1;
var coslack/2;
dicted label for unlabeled instances
#th
var weight/1;
#th
var b/0;
6 var r/O;
#th
Logically parameterized LP variable
    (set of ground LP variables)
#margin
slack = sum{label(I)} slack(I);
coslack = sum{cite(I1,I2),label(I1),query(I2)} slack(I1,I2)
    + sum{cite(I1,I2),label(I2),query(I1)} slack(I1,I2)
#find the largest margin. Here the C's encode trade-off parameters
13 minimize: -r + C(1) * slack + C(2) * coslack;
```


## Logically parameterized LP objective

## Logically parameterized LP constraint

```
#examples should be cone correct side of the hyperplane
subject to forall {I in label(I)}:
    label(I)*(innerProd(I) + b) + slack(I) >= r;
#weights are between -1 and 1
subject to forall {J in attribute(_, J)}: -1 <= weight(J) <= 1;
subject to : r >= 0; #the margin is positive
subject to forall {I in label(I)}: slack(I) >= 0; #slacks are positive
```

```
var pred/1; #predicted label for unlabeled instances
var slack/1; #the slacks
var coslack/2; #slack between neighboring instances
var weight/1; #the slope of the hyperplane
var b/O; #the intercept of the hyperplane
var r/O; #margin
slack = sum{label(I)} slack(I);
coslack = sum{cite(I1,I2),label(I1),query(I2)} slack(I1,I2)
    + sum{cite(I1,I2),label(I2),query(I1)} slack(I1,I2)
#find the largest margin. Here the C's encode trade-off parameters
minimize: -r + C(1) * slack + C(2) * coslack;
```

```
subject to forall {I in query(I)}: pred(I) = innerProd(I) + b;
#related instances should have the same labels
subject to forall {I1, I2 in cite(I1, I2), label(I1), query(I2)}
    label(I1) * pred(I2) + slack(I1, I2) >= r;
subject to forall {I1, I2 in cite(I1, I2), label(I2), query(I1)}:
    label(I2) * pred(I1) + slack(I1, I2) >= r;
#examples shoula be on the correct slae or the nyperpiane
subject to forall {I in label(I)}:
    label(I)*(innerProd(I) + b) + slack(I) >= r;
#weights are between -1 and 1
subject to forall {J in attribute(_, J)}: -1 <= weight(J) <= 1;
subject to : r >= 0; #the margin is positive
subject to forall {I in label(I)}: slack(I) >= 0; #slacks are positive
```

[Kersting, Mladenov, Tokmakov ARXIV 2014, AIJ 2015]

## Cora (most common vs. rest)



## The more observed the more lifting Faster end-to-end even in the light of Gurobi's fast pre-solving heuristics





## Relational MDP LPs



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[Kersting, Mladenov, Tokmakov ARXIV 2014, AIJ 2015]

## Relational MAP LPs

```
var m/2; #single node, pairwise, and
var m/4; #triplewise probabilities
var m/6; #of configurations to be determined by the solver
#value of the MAP assignment
score = sum{w(P, V)} w(P, V) * m(P, V ) +
    sum{w(P1, P2, V1, V2)} w(P1, P2, V1, V2) * m(P1, P2, V1, V2) +
    sum{w(P1, P2, P3, V1, V2, V3)} w(P1, P2, P3, V1, V2, V3) *
    m(P1, P2, P3, V1, V2, V3);
#marginalization of pairwise beliefs
marginalize(P1, P2, V1) = sum{w(P2, V2)} m(P1, P2, V1, V2);
...
#marginalization of ternary beliefs
marginalize(P1, P2, P3, V1) = sum{w(P3, V3), w(P2, V2)}
    m(P1, P2, P3, V1, V2, V3);
...
maximize: score; #find assignment with largest value
subject to forall {P in w(P, _)}:
    sum {w(P, V)} m(P, V) = 1; #atom beliefs sum to one
#pairwise consistency constraints
subject to forall {P1, P2, V1 in w(P1, P2, V1, _)}:
    marginalize(P1, P2, V1) = m(P1, V1);
#ternary consistency constraints
subject to forall {P1, P2, P3, V1 in w(P1, P2, P3, V1, _, _)}:
    marginalize(P1, P2, P3, V1) = m(P1, V1);
```

Stays fix for different relational probabilistic model, which are specified in a knowledge base


## Relational and Compressed Label Propagation

- Many human activities have been shown to exhibit universal patterns.
- What about researcher migration in CS?
- Inferred from 1 million authors of 1,9 million papers


Early Stage: LogNormal


Later Stages:
Gamma


## Brain Circulation: <br> Gamma


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## Relational and Compressed Label Propagation

## Relational way to specify Label Propagation matrix

$$
w_{i j}=w_{i j}+\lambda_{2} \text { if } a(i)=a(j) \wedge y(i)=y(j)
$$

Fractional automorphims for compressing the resulting label propagation matrix; use any label propagation approach (even QP once) on the compressed matrix


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[Kersting, Mladenov, Tokmakov ARXIV 2014, AIJ 2015]

## Relational Mathematical Programming



## Relational Mathematical Programming



## Relational Mathematical Programming



# However, relational programming is not the answer to everything 

Let's embed relational mathematical programming into an imparative programming



# relo $p$ 

http://www-ai.cs.uni-dortmund.de/weblab/static/RLP/htmI/

## RELOOP: A Toolkit for Relational Convex Optimization

| flow = numeric_predicate("flow", 2) cap = numeric_predicate("cap", 2) | Using a probabilistic programming language |
| :---: | :---: |
| model.add_reloop_variable(flow) | stochastic RMPs |
| model += RlpSum([X, Y], source $(X) \&$ | ops and relations get |
| $\begin{aligned} & \text { outFlow }=\operatorname{RlpSum}([X,], \operatorname{edge}(X, Z), f \\ & \text { inFlow }=\operatorname{RlpSum}([Y,], \operatorname{edge}(Z, Y), f] \end{aligned}$ | interwined, and models can refer to each other |
|  |  |

## Take-away Messages

## 1. Graphical models allow to deal with uncertainty

2. Graphs/Matrices are not enough, we need logic /high-level languages
3. Tree-Width is not the end of the story
4. Probabilities are not enough we need optimization
5. Relations and loops should go together

tu

# Democratization of Optimization 

- Reduce the level of expertise necessary to build optimization applications
- Shorten mathematical program code to make models faster to write and easier to communicate
- Reduce development time and cost to encourage experimentation
- Facilitate the construction of more sophisticated models that incorporated rich domain knowledge
- Speed up solvers by exploiting language properties, compression, and compilation

> Achieving this requires the help of all of you! ML, KR, CP, SAT, PL, ... you attention


[^0]:    Kristian Kersting - Democratization of Optimization

