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RELOOP: A Toolkit for Relational Convex Optimization







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Statistical Relational Artificial Intelligence





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Udi

Apsel







and many more ...

The Democratization of Optimization



Kristian Kersting



. . .

Uncertain

Reasoningng

SAT

Optimization

Scaling









Democratization of Data



Take your spreadsheet ...

Features







IS IT REALLY THAT SIMPLE?



NO, e.g., today's data is relational

Features Objects



Nat Rev Genet. 2012 May 2;13(6):395-405

Heart diseases and strokes – cardiovascular disease – are expensive for the world

According to the World Heart Federation, cardiovascular disease cost the European Union EURO169 billion in 2003 and the USA about EURO310.23 billion in direct and indirect annual costs. By comparison, the estimated cost of all cancers is EURO146.19 billion and HIV infections, EURO22.24 billion



Electronic Health Records A New Opportunity for AI to Save our Lifes

[Natarajan, Kersting, et al. IAAI 2013, Springer Briefs in CS 2015, AIME 2015]



Kristian Kersting - Democratization of Optimization

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Statistical Relational AI ...

... the study and design of intelligent agents that act in noisy worlds composed of objects and relations among the objects



Thanks to you - the Iatlian AI community - for your great contributions!

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[Kersting, Driessens ICML 2008; Karwath, Kersting, Landwehr ICDM 2008; Natarajan, Joshi, Tadepelli, Kersting, Shavlik. IJCAI 2011; Khot, Natarajan, Kersting, Shavlik ICDM 2013, MLJ 2012, Springer Brief 2015, MLJ 2015]

Boosted Statistical Relational Learning



Algo	Likelihood	AUC-ROC	AUC-PR	Time
Boosting	0.810	0.961	0.930	9s
MLN	0.730	0.535	0.621	93 hrs

Take-Away Messages

- **1.** Graphical models allow to deal with uncertainty and to make predictions
- 2. Graphs/Matrices are not enough, we need logic/high-level languages



STILL NOT CONVINCED?



Guy van den Broeck's not so simple AI example

What is the probability that the first card of a shuffled deck is an Ace?

Easy for humans but not so easy for graphical models

Exact inference builds a table of ≥13⁵² rows!

Message passing passes ≥ 13⁵² messages!

Graphical model is fully connected, no independencies, high tree-width

Low tree-width is not the final answer

Fast modelling

Fast inference

We need relational nodels w1: $\forall p,x,y, Card(p,x) \land ard(p,y) \Rightarrow x = y$ w2: $\forall c,x,y, Card(x,c) \land Card(y,c) \Rightarrow x = y$ and symmetry-reduction

Take-away Messages

- 1. Graphical models allow to deal with uncertainty
- 2. Graphs/Matrices are not enough, we need logic /high-level languages
- **3.** Tree-Width is not the end of the story



Let's mathematically characterize symmetries for approximate inference

[Singla, Domingos AAAI 2008; Kersting, Ahmadi, Natarajan UAI 2009; Ahmadi, Kersting, Mladenov, Natarajan MLJ 2013]

Lifted Loopy Belief Propagation = Exploit computational symmetries



Compression: Coloring the graph



Color nodes according to the evidence you have

- No evidence, say red
- State "one", say brown
- State "two", say orange
- • •

Color factors distinctively according to their equivalence classes. For instance, assuming f_1 and f_2 to be identical and B appears at the second position within both, say blue



[Singla, Domingos AAAI 2008; Kersting, Ahmadi, Natarajan UAI 2009; Ahmadi, Kersting, Mladenov, Natarajan MLJ 2013]

Compression: Pass the colors around



1. Each factor collects the colors of its neighboring nodes





1. Each factor collects the colors of its neighboring nodes

2. Each factor "signs" ist color signature with its own color





- 1. Each factor collects the colors of its neighboring nodes
- 2. Each factor "signs" ist color signature with its own color
- 3. Each node collects the signatures of its neighboring factors





- 1. Each factor collects the colors of its neighboring nodes
- 2. Each factor "signs" ist color signature with its own color
- 3. Each node collects the signatures of its neighboring factors
- 4. Nodes are recolored according to the collected signatures



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- 1. Each factor collects the colors of its neighboring nodes
- 2. Each factor "signs" ist color signature with its own color
- 3. Each node collects the signatures of its neighboring factors
- 4. Nodes are recolored according to the collected signatures
- 5. If no new color is created stop, otherwise go back to 1 Kristian Kersting - Democratization of Optimization

Compression: ... and compute the quotient factor graph



Essentially we just compute the so-called quotient factor graph

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[Singla, Domingos AAAI 2008; Kersting, Ahmadi, Natarajan UAI 2009; Ahmadi, Kersting, Mladenov, Natarajan MLJ 2013]

Finally, run a modified Loopy Belief Propagation



- Nodes are now groups of random variables
- The counts ensure that we send the same number of message as standard loopy belief propagation



Lifted Probabilistic Inference



It turns out that color passing is well known in graph theory

The Weisfeiler Lehman Algorithm

Weisfeiler-Lehman (WL) Algorithmus aka "naive vertex classification"

- Basic subroutine for GI testing
- Computes LP-relaxations of GA-ILP, aka. fractional automorphisms
- Quasi-linear running time O((n+m)log(n)) when using asynchronous updates [Berkholz, Bonsma, Grohe ESA 2013]



- Part of graph tool SAUCY [See e.g. Darga, Sakallah, Markov DAC 2008]
- Can be extended to weighted graphs/real-valued matrices [Grohe, Kersting, Mladenov, Selman ESA 2014]
- Actually a Frank-Wolfe optimizer and can be viewed as recursive spectral clustering
 [Kersting, Mladenov, Garnett, Grohe AAAI 2014]





From Factor Graphs to Graphs



Encode the factor colors into the node colors



Then run Weisfeiler Lehman / Color Passing just on the graph



Instead of looking at AI through the glasses of probabilities over possible

probabilities over possible worlds, we may also approach it using optimization [Mladenov, Ahmadi, Kersting AISTATS 2012, Grohe, Kersting, Mladenov, Selman ESA 2014, Kersting Mladenov, Tokmatov AIJ 2015]

Compressing Linear Programs



(1) Reduce the LP by running WL on the LP-Graph(2) Run any solver on the (hopefully) smaller LP



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quasi-linear overhead that may result in exponential speed up

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As also noted by Stephen Boyd

DENSE VS. SPARSE IS NOT ENOUGH, SOLVERS NEED TO BE AWARE OF SYMMETRIES





[Mladenov, Ahmadi, Kersting AISTATS 2012, Grohe, Kersting, Mladenov, Selman ESA 2014, Kersting Mladenov, Tokmatov AIJ 2015]

Compute Equitable Partition (EO) of the LP using WL



$$\mathcal{P} = \{P_1, \dots, P_p; Q_1, \dots, Q_q\}$$

Partition of Partition of LP variables LP constraints

Intuitively, we group together variables resp. constraints that interact in the very same way in the LP.

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technische universität dortmund **Fractional Automorphisms of LPs**

The EP induces a fractional automorphism of the coefficient matrix ${\bf A}$

$$\mathbf{X}_{\boldsymbol{Q}}\mathbf{A} = \mathbf{A}\mathbf{X}_{\boldsymbol{P}}$$

where X_Q and X_p are doubly-stochastic matrixes (relaxed form of automorphism)

$$\begin{aligned} (\mathbf{X}_{P})_{ij} &= \begin{cases} 1/|P| & \text{if both vertices } i, \ j \text{ are in the same } P, \\ 0 & \text{otherwise.} \end{cases} \\ (\mathbf{X}_{Q})_{ij} &= \begin{cases} 1/|Q| & \text{if both vertices } i, \ j \text{ are in the same } Q, \\ 0 & \text{otherwise} \end{cases} \end{aligned}$$

Fractional Automorphisms Preserve Solutions

If **x** is feasible, then $\mathbf{X}_{p}\mathbf{x}$ is feasible, too.

By induction, one can show that left-multiplying with a double-stochastic matrix preserves directions of inequalities. Hence,

$\mathbf{A}\mathbf{x} \leq \mathbf{b} \Rightarrow \mathbf{X}_{\boldsymbol{Q}}\mathbf{A}\mathbf{x} \leq \mathbf{X}_{\boldsymbol{Q}}\mathbf{b} \Leftrightarrow \mathbf{A}\mathbf{X}_{\boldsymbol{P}}\mathbf{x} \leq \mathbf{b}$



Fractional Automorphisms Preserve Solutions

If \mathbf{x}^* is optimal, then $\mathbf{X}_p \mathbf{x}^*$ is optimal, too. Since by construction $\mathbf{c}^T \mathbf{X}_P = \mathbf{c}^T$ and hence $\mathbf{c}^T (\mathbf{X}_P \mathbf{x}) = \mathbf{c}^T \mathbf{x}$



What have we established so far?

Instead of considering the original LP $(\mathbf{A},\mathbf{b},\mathbf{c})$

It is sufficient to consider

 $(\mathbf{A}\mathbf{X}_{P}, \mathbf{b}, \mathbf{X}_{P}{}^{T}\mathbf{c})$

i.e. we "average" parts of the polytope.

But why is this dimensionality reduction?



Dimensionality Reduction

The doubly-stochastic matrix \mathbf{X}_P can be written

as
$$\mathbf{X}_P = \mathbf{B}\mathbf{B}^T$$

 $\mathbf{B}_{iP} = \begin{cases} \frac{1}{\sqrt{|P|}} & \text{if vertex } i \text{ belongs to part } P \\ 0 & \text{otherwise.} \end{cases}$

Since the column space of B is equivalent to the span of \mathbf{X}_P , it is actually sufficient to consider only $(\mathbf{AB}_P, \mathbf{b}, \mathbf{B}_P^T \mathbf{c})$

This is of reduced size and actually we can also drop any constraint that becomes identical

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Dimensionality Reduction of LPs



Feasible region of LP and the objective vectors

Span of the fractional automorpishm of the LP Projections of the feasible region onto the span of the fractional automorphism



Fractional automorphisms provide an algebraic tool to study lifted messagepassing approaches

 $= \mathbf{A} \mathbf{X}_{P}$

Any MAP-LP message-passing approach is liftable





[Mladenov, Kersting UAI 2015]

Any concave free energy is liftable 8000 읍 7500-) 7000-) 6500-10 around <u>ල</u> 3.0 2.5 ti 2.0 |V| + |F|, log •••• reparam 8 Opjective 6000 5500 5500 Running 1.0 0.5 0.0 450<u>0</u> 30 -20 -10 100 150 200 250 300 20 50 100 150 200 250 300 0 W 10 30 50 Domain Size Domain Size (a) Complete Graph MLN. 63700r ຄົ 4.0 63600 ground 12 3.5 3.0 2.5 2.5 2.0 1.5 1.5 1.0 ਦੋਂ 63500 •••• reparam ් ₆₃₄₀₀ e 63300 e 63200 6 63100 63200 6300p 、0.5∟ 0.5 1.5 2.0 茍 50 100 150 200 250 300 50 100 150 200 250 300 1.0 Domain Size Domain Size (b) Clique-Cycle MLN. 140000 scale 14 <u>ല</u>,120000 around 4 12 <u>මි</u> 10 log ਚੋ100000 •••• reparam 00008 ق time, Ę, Objective 60000 + nning 40000 2000 $\frac{0}{10}$ 100 150 200 250 300 茍 50 100 150 200 250 300 -5 10 6 50 0 W 5

Actually, this is the first distributed lifted message-passing approach

Domain Size

(c) Friends-smokers MLN.

Domain Size

All MP Inference Approaches are Liftable





However, it is annoying that we still have to develop modified message passing approaches tation

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enley

Fractional automorphisms allow one to eliminate this reliance [Mladenov, Globerson, Kersting UAI 2014; Mladenov, Kersting UAI 2015]







[Bennett 1999; Mangasarian 1999; Zhou, Zhang, Jiao 2002, ...]

Classification using LPs



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Replace I_2 - by I_1 -, I_{∞} -norm in the standard SVM formulation

[Kersting, Mladenov, Tokmakov ARXIV 2014, AIJ 2015]

Declarative Machine Learning



[Kersting, Mladenov, Tokmakov ARXIV 2014, AIJ 2015]

Declarative Machine Learning 2.0

```
1 var pred/1;
                          #predicted label for unlabeled instances
 Abstract LP-SVM
      2 var slack/1:
                          #the slacks
      3 var coslack/2;
                          #slack between neighboring instances
      4 var weight/1;
                          #the slope of the hyperplane
      5 var b/0;
                          #the intercept of the hyperplane
      6 \text{ var } r/0;
                          #margin
      \overline{7}
      8 slack = sum{label(I)} slack(I);
      9 coslack = sum{cite(I1,I2),label(I1),query(I2)} slack(I1,I2)
                + sum{cite(I1,I2),label(I2),guery(I1)} slack(I1,I2)
      10
      11
     12 #find the largest margin. Here the C's encode trade-off parameters
     13 minimize: -r + C(1) + slack + C(2) + coslack;
     15 subject to forall {I in guery(I)}: pred(I) = innerProd(I) + b;
Collective
     16 #related instances should have the same labels.
     17 subject to forall {I1, I2 in cite(I1, I2), label(I1), query(I2)}
     18 label(I1) * pred(I2) + slack(I1, I2) >= r;
                                                                     relational
     19 #the symmetric case
     20 subject to forall {I1, I2 in cite(I1, I2), label(I2), query(I1)}:
  Ô
          label(I2) * pred(I1) + slack(I1, I2) >= r;
     21
      22
      23 #examples should be on the correct side of the hyperplane
      24 subject to forall {I in label(I)}:
            label(I)*(innerProd(I) + b) + slack(I) >= r;
      25
      26 #weights are between -1 and 1
      27 subject to forall {J in attribute(_, J)}: -1 <= weight(J) <= 1;</pre>
     28 subject to : r \ge 0;
                                     #the margin is positive
      29 subject to forall {I in label(I)}: slack(I) >= 0; #slacks are positive
```

Cora (most common vs. rest)



The more observed the more lifting Faster end-to-end even in the light of Gurobi's fast pre-solving heuristics



U technische universität dortmund [Kersting, Mladenov, Tokmakov ARXIV 2014, AIJ 2015]

Relational MDP LPs



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[Kersting, Mladenov, Tokmakov ARXIV 2014, AIJ 2015] Relational MAP LPS

```
1 var m/2; #single node, pairwise, and
2 var m/4; #triplewise probabilities
3 \text{ var m/6}; #of configurations to be determined by the solver
4 #value of the MAP assignment
5 score = sum\{w(P, V)\}w(P, V) * m(P, V) +
    sum{w(P1, P2, V1, V2)} w(P1, P2, V1, V2) * m(P1, P2, V1, V2) +
6
   sum{w(P1, P2, P3, V1, V2, V3)} w(P1, P2, P3, V1, V2, V3) *
7
    m(P1, P2, P3, V1, V2, V3);
8
9
10 #marginalization of pairwise beliefs
11 marginalize(P1, P2, V1) = sum\{w(P2, V2)\}m(P1, P2, V1, V2);
12 . . .
13 #marginalization of ternary beliefs
14 marginalize(P1, P2, P3, V1) = sum\{w(P3, V3), w(P2, V2)\}
15 m(P1, P2, P3, V1, V2, V3);
16 . . .
                                             #find assignment with largest value
17 maximize: score:
18 subject to forall {P in w(P, _)}:
      sum \{w(P, V)\} m(P, V) = 1;
                                             #atom beliefs sum to one
19
20 #pairwise consistency constraints
21 subject to forall {P1, P2, V1 in w(P1, P2, V1, _)}:
      marginalize(P1, P2, V1) = m(P1, V1);
22
23 . . .
24 #ternary consistency constraints
25 subject to forall {P1, P2, P3, V1 in w(P1, P2, P3, V1, _, _)}:
      marginalize(P1, P2, P3, V1) = m(P1, V1);
26
```

Stays fix for different relational probabilistic model, which are specified in a knowledge base

What about relational QPs, SDPs, ...? What about lifted solvers for them? What about knowledge compilation for optimization?

Of course, the focus is not just linear programs!

Relational and Compressed Label Propagation

- Many human activities have been shown to exhibit universal patterns.
- What about researcher migration in CS?
- Inferred from 1 million authors of 1,9 million papers



Relational and Compressed Label Propagation

Relational way to specify Label Propagation matrix

$$w_{ij} = w_{ij} + \lambda_2$$
 if a (i) = a (j) \wedge y (i) = y (j)

Fractional automorphims for compressing the resulting label propagation matrix; use any label propagation approach (even QP once) on the compressed matrix





Relational Mathematical Programming





Relational Mathematical Programming





Relational Mathematical Programming



However, relational programming is not the answer to everything

Let's embed relational mathematical programming into an imparative programming



http://www-ai.cs.uni-dortmund.de/weblab/static/RLP/html/ RELOOP: A Toolkit for Relational Convex Optimization

model = RlpProblem("flow", LpMaximize, PvDatalogLogKb(), Pulp) Using a probabilistic flow = numeric_predicate("flow", 2) programming language cap = numeric_predicate("cap", 2) we can even get model.add_reloop_variable(flow) stochastic RMPs model += RlpSum([X, Y], source(X) & e Loops and relations get interwined, and models outFlow = RlpSum([X,], edge(X, Z), inFlow = RlpSum([Y,], edge(Z, Y), f]can refer to each other model += ForAll([Z,], node(Z) & \sim source(Z) & \sim target(Z), inFlow |eq| outFlow)

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Take-away Messages

- 1. Graphical models allow to deal with uncertainty
- 2. Graphs/Matrices are not enough, we need logic /high-level languages
- **3.** Tree-Width is not the end of the story
- **4.** Probabilities are not enough we need optimization
- **5.** Relations and loops should go together



Democratization of Optimization

- Reduce the level of expertise necessary to build optimization applications
- Shorten mathematical program code to make models faster to write and easier to communicate
- Reduce development time and cost to encourage experimentation
- Facilitate the construction of more sophisticated models that incorporated rich domain knowledge
- Speed up solvers by exploiting language properties, compression, and compilation

Achieving this requires the help of [Thanks for all of you! ML, KR, CP, SAT, PL, ... you attention

